

Week 3: Lab

COLLABORATION LEVEL 0 (NO RESTRICTIONS). OPEN NOTES.

1. Consider the linear-time merge algorithm discussed in the notes and a possible implementation below. How many element comparisons will the standard merge function take to merge the following left and right lists ?

$left = [1, 3, 4, 5, 6, 7, 8]$, $right = [1, 5, 9, 11, 12, 16]$

Merge(left, right)

```
result = []
i=0
j=0
while i < len(left) and j < len(right):
    if left[i] < right[j]:
        result.append(left[i])
        i = i+1
    else:
        result.append(right[j])
        j = j+1

# add any left overs
while i < len(left):
    result.append(left[i])
    i = i+1
while j < len(right):
    result.append(right[j])
    j = j+1

return result
```

- A 8
- B 9
- C 10
- D 13

Find a $\Theta()$ bound for the following recurrences using iteration. Assume $T(1) = 1$.

*What we expect: show $O(1)$ steps of your iteration, with the general formula, the recursion depth, and the final $\Theta()$ bound for $T(n)$. **Do not write your answers on this page.** Use a separate page for each problem and show your work.*

2. $T(n) = T(n/2) + 1$

3. $T(n) = T(n/3) + 1$ (assume $T(i) = 1$ for $i = 1, 2$).

4. $T(n) = T(n/10) + 1$ (assume $T(i) = 1$ for $i < 10$).

5. $T(n) = T(2n/3) + 1$ (assume $T(i) = 1$ for $i = 1, 2$).

6. $T(n) = T(n - 1) + 1$

7. $T(n) = T(n - 2) + 1$ (assume $T(i) = 1$ for $i = 1, 2$).

8. $T(n) = T(n - 3) + 1$ (assume $T(i) = 1$ for $i = 1, 2, 3$).

9. $T(n) = T(n/2) + n$

10. $T(n) = T(n/3) + n$ (assume $T(i) = 1$ for $i < 3$).

11. $T(n) = 3T(n/3) + \Theta(n)$ (assume $T(i) = 1$ for $i < 3$).

12. $T(n) = 5T(n/5) + \Theta(n)$ (assume $T(i) = 1$ for $i < 5$).

13. $T(n) = T(n - 1) + n$

14. $T(n) = T(n - 2) + \Theta(n)$ (assume $T(i) = 1$ for $i = 1, 2$).

15. $T(n) = T(n - 1) + 2n - 3$, with $(T(1) = 1)$

16. $T(n) = T(\sqrt{n}) + 1$
17. $T(n) = 7T(n/2) + n^3$
18. $T(n) = 7T(n/2) + n^2$
19. $T(n) = 4T(n/3) + 2n - 1$, with $(T(1) = T(2) = 1$
20. $T(n) = 3T(n/2) + n^2$, with $(T(1) = 1$
21. $T(n) = 2T(n - 1) + \Theta(1)$
22. (challenge) $T(n) = T(n/3) + T(2n/3) + \Theta(n)$ (only guess the solution)
23. (challenge) $T(n) = T(n/3) + T(n/4) + \Theta(n)$ (only guess the solution)
24. (challenge) $T(n) = T(n/2) + T(n/4) + T(n/10) + \Theta(n)$ (only guess the solution)
25. Based on all examples seen so far, list recurrences that solve to:
 - (a) $\Theta(\lg n)$
 - (b) $\Theta(n)$
 - (c) $\Theta(n \lg n)$
 - (d) $\Theta(n^2)$
 - (e) exponential

For each category, enumerate all recurrences seen so far that fall into that category, and add at last one *new* one.

Partial Answers

1. B (9 comparisons)
2. $T(n) = T(n/2) + 1 : \Theta(\lg n)$
3. $T(n) = T(n/3) + 1$ (assume $T(i) = 1$ for $i = 1, 2$): $\Theta(\lg n)$
4. $T(n) = T(n/10) + 1$ (assume $T(i) = 1$ for $i < 10$): $\Theta(\lg n)$
5. $T(n) = T(2n/3) + 1$ (assume $T(i) = 1$ for $i = 1, 2$): $\Theta(\lg n)$
6. $T(n) = T(n - 1) + 1 : \Theta(n)$
7. $T(n) = T(n - 2) + 1$ (assume $T(i) = 1$ for $i = 1, 2$): $\Theta(n)$
8. $T(n) = T(n - 3) + 1$ (assume $T(i) = 1$ for $i = 1, 2, 3$): $\Theta(n)$
9. $T(n) = T(n/2) + n : \Theta(n)$
10. $T(n) = T(n/3) + n$ (assume $T(i) = 1$ for $i < 3$): $\Theta(n)$
11. $T(n) = 3T(n/3) + \Theta(n)$ (assume $T(i) = 1$ for $i < 3$): $\Theta(n \lg n)$
12. $T(n) = 5T(n/5) + \Theta(n)$ (assume $T(i) = 1$ for $i < 5$): $\Theta(n \lg n)$
13. $T(n) = T(n - 1) + n : \Theta(n^2)$
14. $T(n) = T(n - 2) + \Theta(n)$ (assume $T(i) = 1$ for $i = 1, 2$): $\Theta(n^2)$
15. $T(n) = T(n - 1) + 2n - 3$, with $(T(1) = 1) : \Theta(n^2)$
16. $T(n) = T(\sqrt{n}) + 1 : \Theta(\lg \lg n)$
17. $T(n) = 7T(n/2) + n^3 : \Theta(n^3)$
18. $T(n) = 7T(n/2) + n^2 : T(n) = \Theta(n^{\log_7 8})$
19. $T(n) = 4T(n/3) + 2n - 1$, with $(T(1) = T(2) = 1) : T(n) = \Theta(n^{\log_3 4})$
20. $T(n) = 3T(n/2) + n^2$, with $(T(1) = 1) : T(n) = \Theta(n^2)$
21. $T(n) = 2T(n - 1) + \Theta(1) : T(n) = \Theta(2^n)$ Note: For exponential recurrences we are usually happy with just a lower bound.
22. (challenge) $T(n) = T(n/3) + T(2n/3) + \Theta(n) : n/a$
23. (challenge) $T(n) = T(n/3) + T(n/4) + \Theta(n) : n/a$
24. (challenge) $T(n) = T(n/2) + T(n/4) + T(n/10) + \Theta(n) : n/a$